

Contact topology and holomorphic invariants via elementary combinatorics

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16 April 2013

Outline

- 1 Introduction
- 2 Combinatorial and algebraic structure
- 3 Contact topology
- 4 Holomorphic invariants

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- 1 Introduction
 - Overview
 - Symplectic geometry
 - Contact geometry
 - Complex structures and holomorphic curves
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Overview

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 - moduli spaces of pseudo-holomorphic curves
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 - intricate algebraic structures keeping track of analytic data

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- Much of it is quite involved, requiring:
 - Fredholm / index theory of Cauchy-Riemann operators
 - moduli spaces of pseudo-holomorphic curves
 - delicate differential geometry and topology
 - intricate algebraic structures keeping track of analytic data
- However, *in the simplest cases* some of this structure reduces to some *elementary combinatorics and algebra* which is interesting in its own right.

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- Give some *very* brief background to the subjects of symplectic and contact geometry and holomorphic curves.
- Discuss some of our algebraic and combinatorial results in their own right.
(No symplectic/contact geometry or holomorphic curves assumed.)
- Briefly explain how this elementary algebra/combinatorics describes contact topology and arises from holomorphic invariants.

Symplectic manifolds

Definition

A symplectic manifold is a pair

$$(M, \omega)$$

where

- *M is a smooth manifold*
- *ω is a closed 2-form ($d\omega = 0$) which is non-degenerate.*

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Structure of Hamiltonian mechanics:

- Given a smooth function $H : M \rightarrow \mathbb{R}$ (Hamiltonian) we obtain a 1-form dH and a dual vector field X_H via

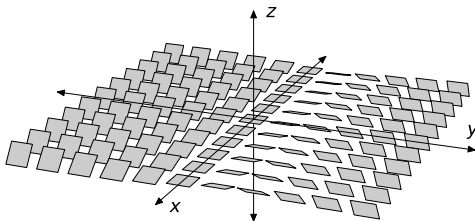
$$\omega(X_H, \cdot) = dH$$

E.g. $M = \mathbb{R}^{2n}$, $\omega = \sum_{j=1}^n dx_j \wedge dy_j$.

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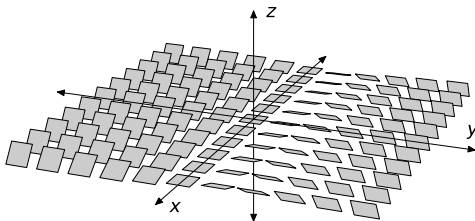


Figure 3.3 illustrates the three types of functions that can be used to model data. The first type is a linear function, which is a straight line. The second type is a quadratic function, which is a parabola. The third type is an exponential function, which is a curve that increases or decreases rapidly.

E.g. \mathbb{R}^3 with $\alpha = dz - y \, dx$.

Symplectic vs complex geometry

- Complex geometry also only exists in *even* number of dimensions.
- Gromov (1985): Consider *almost complex structures* on symplectic manifolds and *holomorphic curves*.

Symplectic vs complex geometry

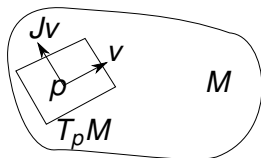
- Complex geometry also only exists in *even* number of dimensions.
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Definition

An almost complex structure on a smooth manifold is a map

$$J : TM \longrightarrow TM$$

preserving each fibre $T_p M$ and satisfying $J^2 = -1$.



Almost complex vs complex

- *Almost complex structure* is a pointwise definition.
- A *complex structure* requires local charts to \mathbb{C}^n with holomorphic transition maps.
(Much more onerous.)

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- Every symplectic manifold has a *compatible* almost complex structure J , and all choices of compatible J are homotopic.

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Existence:

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- Every symplectic manifold has a *compatible* almost complex structure J , and all choices of compatible J are homotopic.
(Compatible: J and ω behave in linear algebra like i and $dx \wedge dy$. $\omega(v, w) = \omega(Jv, Jw)$ and $\omega(v, Jv) > 0$)

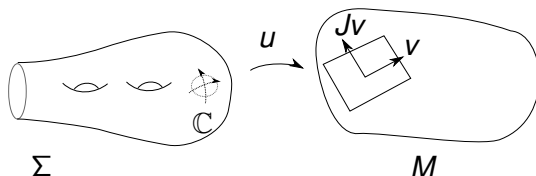
Holomorphic curves

Given symplectic (M, ω) and compatible almost complex $J...$

Definition

A holomorphic curve is a map $u : \Sigma \rightarrow M$, where Σ is a Riemann surface, satisfying the Cauchy-Riemann equations

$$Du \circ j = J \circ Du.$$



An *almost* complex structure is sufficient for the equations:
“pseudo-holomorphic”, “ J -holomorphic”.

Moduli spaces

- Given appropriate constraints (marked points, boundary conditions) and transversality, the space of holomorphic curves is a finite-dimensional orbifold: *moduli space* \mathcal{M} .

Moduli spaces

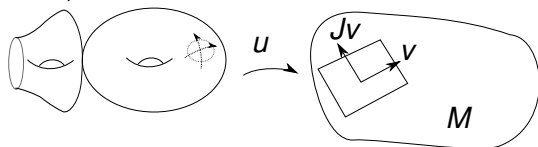
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- \mathcal{M} compactified to $\overline{\mathcal{M}}$: *Gromov compactness theorem*.
- Boundary of $\overline{\mathcal{M}}$ is stratified: boundary strata are moduli spaces for “degenerate” holomorphic curves (nodal surfaces, etc.)



Homology theories

Floer Homology theories (e.g. contact homology, Heegaard Floer homology), roughly...

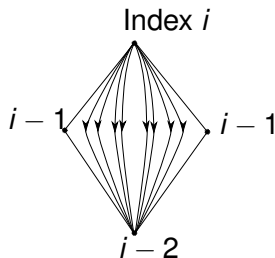
- Define a chain complex generated by boundary conditions for holomorphic curves
- A differential counting 0-dimensional families of holomorphic curves between boundary conditions.
- Boundary structure of moduli space gives $\partial^2 = 0$.

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(Analogous Morse construction of singular homology: complex generated by critical points of Morse function, differential counts 0-dimensional families of gradient trajectories.)



The power of holomorphic invariants

- Floer homology theories give very powerful invariants of 3-manifolds, knots, etc...
- Related to Seiberg–Witten theory, Donaldson–Thomas theory, etc...
- E.g., *knot Floer homology* can compute the genus of a knot.

The power of holomorphic invariants

- Floer homology theories give very powerful invariants of 3-manifolds, knots, etc...
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- E.g., *knot Floer homology* can compute the genus of a knot.
- For a *less complicated* variant called *sutured Floer homology*, and a *simple class* of manifolds $M = \Sigma \times S^1$, we obtain all the combinatorial structure we are about to see, and more...

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 - Quantum Pawn Dynamics (QPD)
 - Adjoining adjoints
 - Chord diagrams
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Quantum Pawn Dynamics

- Pawns on a finite 1-dimensional chessboard.
- A state of the QPD universe:



- Pawns move from left to right, one square at a time.
(No capturing, no en passant, no double first moves.)

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Quantum pawns: “Inner product” $\langle \cdot | \cdot \rangle$ describes the possibility of pawn moves from one state to another.
Valued in \mathbb{Z}_2 .

Definition (Pawn “inner product”)

$$\langle w_0 | w_1 \rangle = \begin{cases} 1 & \text{if it is possible for pawns to move from } w_0 \text{ to } w_1 \\ & \text{(this includes the case } w_0 = w_1 \text{);} \\ 0 & \text{if not.} \end{cases}$$

E.g.

$$\langle \begin{array}{|c|c|c|c|c|} \hline \text{♙} & & \text{♙} & \text{♙} & \\ \hline \end{array} \mid \begin{array}{|c|c|c|c|c|} \hline & \text{♙} & \text{♙} & & \text{♙} \\ \hline \end{array} \rangle = 1$$

Quantum Pawn Dynamics

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Also, entangled chessboards.

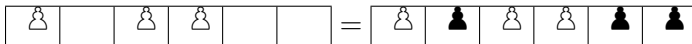
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Note asymmetry of $\langle \cdot | \cdot \rangle$.

A “booleanized” partial order. (Complete lattice.)

Dirac Pawn Sea

- Think of an “empty” chessboard as a thriving sea of anti-pawns.
“Anti-pawn” = “absence of pawn”.



Creation and annihilation operators

The *initial pawn creation operator* $a_{p,0}^*$ adjoins a new *initial* (leftmost) square to the chessboard, containing a pawn.

$$a_{p,0}^* \begin{array}{|c|c|c|c|c|c|} \hline \text{white} & \text{black} & \text{white} & \text{white} & \text{black} & \text{black} \\ \hline \end{array} = \begin{array}{|c|c|c|c|c|c|c|} \hline \text{white} & \text{white} & \text{black} & \text{white} & \text{white} & \text{black} & \text{black} \\ \hline \end{array}$$

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The *initial pawn annihilation operator* $a_{p,0}$ deletes the leftmost square from the chessboard, and a pawn on it.

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If no pawn (anti-pawn) in the leftmost square, try to delete...
“error 404 universe not found” mod 2 = 0.

$$a_{p,0} \begin{array}{|c|c|c|c|c|} \hline \text{black} & \text{white} & \text{white} & \text{black} & \text{black} \\ \hline \end{array} = 0$$

Similar initial anti-pawn annihilation $a_{q,0}$ and creation $a_{q,0}^\dagger$.

Creation of chessboards







- The *vacuum* state of the QPD universe is the null chessboard \emptyset .
(Note $\emptyset \neq 0$.)
- Applying initial creation operators to the vacuum can create any chessboard.

$$a_{p,0}^* a_{q,0}^\dagger a_{p,0}^* a_{p,0}^* a_{q,0}^\dagger a_{q,0}^\dagger \quad \emptyset \quad = \quad \begin{array}{|c|c|c|c|c|c|} \hline \text{white king} & \text{black king} & \text{white king} & \text{white king} & \text{black king} & \text{black king} \\ \hline \end{array}$$

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- The $*$ and † refer to *adjoints*.
(*Galois connections* on partial orders.)

Adjoint

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- As our “inner product” is asymmetric, we have *two distinct adjoints* f^*, f^\dagger of an operator f .

$$\langle fx|y\rangle = \langle x|f^*y\rangle, \quad \langle x|fy\rangle = \langle f^\dagger x|y\rangle.$$

So $f^{*\dagger} = f^{\dagger*} = f$.

Initial creation and annihilation are adjoint

Proposition

$$\langle a_{p,0}x|y\rangle = \langle x|a_{p,0}^*y\rangle$$

Proof.

$a_{p,0}^*y$ begins with a pawn.

If x begins with an anti-pawn, both sides are 0.

If x begins with a pawn, $\langle x|a_{p,0}^*y\rangle \neq 0$ compares two chessboards with initial pawns.

$a_{p,0}$ removes an initial pawn so $\langle a_{p,0}x|y\rangle$ gives the same result. □

Similarly, initial anti-pawn creation/annihilation \dagger -adjoint.

Adjoining adjoints

What is $a_{p,0}^{**}$? What operator f satisfies

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 &= \langle \quad \boxed{} \boxed{\text{p}} \boxed{\text{p}} \boxed{} \boxed{} \quad | \quad \boxed{} \boxed{\text{p}} \boxed{} \boxed{\text{p}} \boxed{} \quad \rangle = 1
 \end{aligned}$$

Iterated adjoints

Proposition

The iterated adjoints of $a_{p,0}$ are

$$a_{p,0} \rightarrow a_{p,0}^* \rightarrow a_{p,1} \rightarrow a_{p,1}^* \rightarrow a_{p,2} \rightarrow \cdots \rightarrow a_{p,\Omega} \rightarrow a_{p,\Omega}^*$$

where:

$a_{p,i}$ deletes the i 'th pawn

$a_{p,i}^$ doubles the i 'th pawn*

$a_{p,\Omega}, a_{p,\Omega}^$ are final pawn creation and annihilation.*

Similarly for anti-pawns in the opposite direction.

$$a_{q,\Omega}^\dagger \rightarrow a_{q,\Omega} \rightarrow \cdots a_{q,2} \rightarrow a_{q,1}^\dagger \rightarrow a_{q,1} \rightarrow a_{q,0}^\dagger \rightarrow a_{q,0}$$

(A simplicial structure.)

Adjoint periodicity

Hence

$$a_{p,0}^{*2n_p+2} = a_{p,\Omega}$$

where n_p = number of pawns.

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Theorem (M.)

$$a_{p,0}^{*2n+2} = a_{p,0}.$$

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Adjoint periodicity

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Theorem (M.)

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One can also show that the *duality* operator defined by

$$\langle u|v\rangle = \langle v|Hu\rangle$$

satisfies

Theorem (M.)

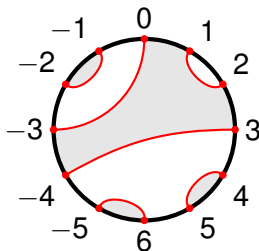
$$H^{2n+2} = 1.$$

Chord diagrams

Consider a disc D with some points F marked on ∂D .

A *chord diagram* is a collection of non-intersecting curves on D joining points of F .

E.g.

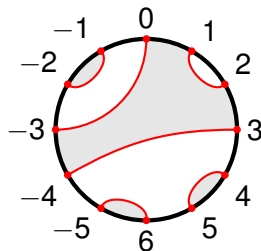


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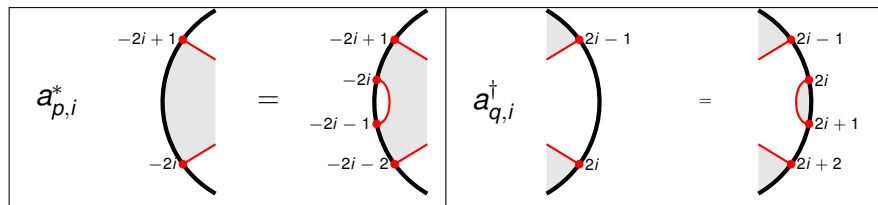
E.g.



- Curves join points of opposite parity, so shade as shown.
- 0 is a basepoint.

Creation and annihilation of chords

Define *creation operators* $a_{p,i}^*$, $a_{q,i}^\dagger$ to insert a new chord in a specific place in a chord diagram as shown.

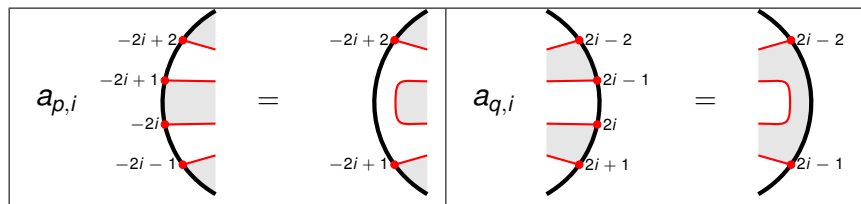


$a_{p,i}^*$ creates a *white* region i spots down on the left.

$a_{q,i}^\dagger$ creates a *black* region i spots down on the right.

Creation and annihilation of chords

Define *annihilation operators* $a_{p,i}$, $a_{q,i}$ to close off chords in a chord diagram as shown.



$a_{p,i}$ closes off a black region i spots down on the left.

$a_{q,i}$ closes off a white region i spots down on the right.

Diagrams of chessboards

The simplest chord diagram is called the *vacuum* Γ_\emptyset .



Build up more complicated diagrams with creation operators.

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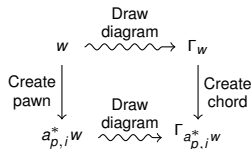


Build up more complicated diagrams with creation operators.

Proposition (M.)

For any chessboard w , there is a chord diagram Γ_w such that creation and annihilation operators agree (are equivariant):

$$\Gamma_{a_{p,i}^* w} = a_{p,i}^* \Gamma_w.$$



Ski slopes

Construction of the *slalom skiing* chord diagram of a chessboard.

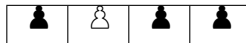


Ski slopes

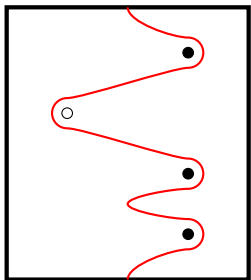
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qpqq

\leftrightarrow



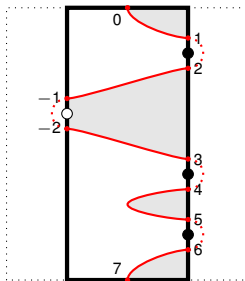
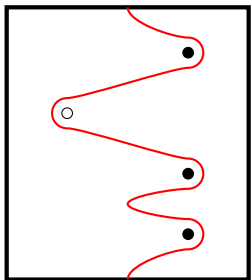
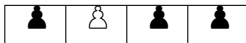
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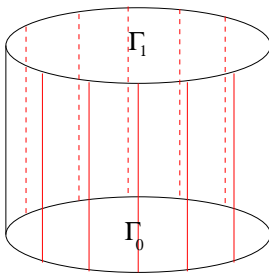
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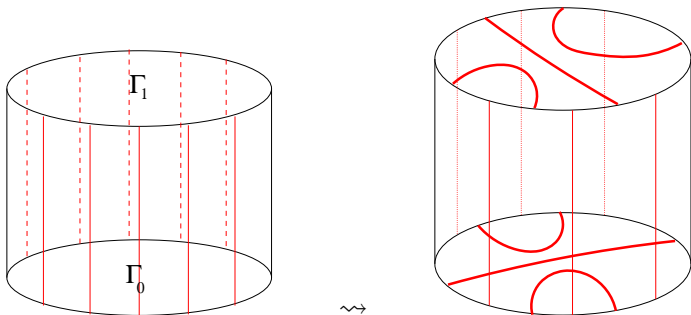
An “Inner product” on chord diagrams

There's a bilinear form on chord diagrams defined by *entering into a cylinder*.



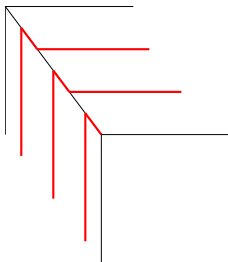
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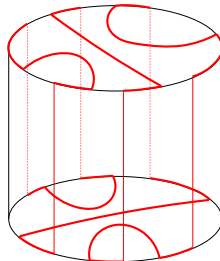
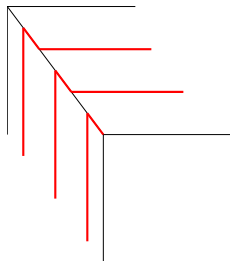
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Definition

$$\langle \Gamma_0 | \Gamma_1 \rangle = \begin{cases} 1 & \text{if the resulting curves on the cylinder} \\ & \text{form a single connected curve;} \\ 0 & \text{if the result is disconnected.} \end{cases}$$

Theorem (M.)

For any two chessboards w_0, w_1 ,

$$\langle w_0 | w_1 \rangle = \langle \Gamma_{w_0} | \Gamma_{w_1} \rangle.$$

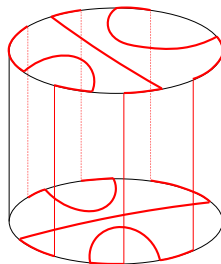
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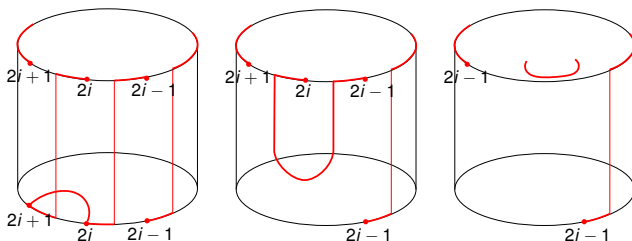
E.g.

$$\left\langle \begin{array}{|c|c|} \hline \text{white} & \text{black} \\ \hline \end{array} \mid \begin{array}{|c|c|} \hline \text{black} & \text{white} \\ \hline \end{array} \right\rangle = 1 =$$



Adjoint

Adjoint relations can be seen *topologically* as “finger moves”.



$$\langle a_{q,i}^\dagger \Gamma_0 \mid \Gamma_1 \rangle = \langle \Gamma_0 \mid a_{q,i} \Gamma_1 \rangle$$

Now perhaps believable that adjoint is periodic.

Bypass surgery

In a chord diagram on disc D , consider a sub-disc B as shown:



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Two natural ways to adjust this chord diagram, consistent with the colours: *bypass surgeries*.

 Γ'  Γ  Γ''

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 Γ'

 Γ

 Γ''

Proposition

With $\Gamma, \Gamma', \Gamma''$ as above, for any Γ_1 ,

$$\langle \Gamma | \Gamma_1 \rangle + \langle \Gamma' | \Gamma_1 \rangle + \langle \Gamma'' | \Gamma_1 \rangle = 0.$$

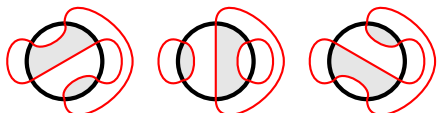
Bypass surgery

Idea of proof:

$$1 + 0 + 1 = 0$$

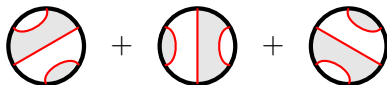
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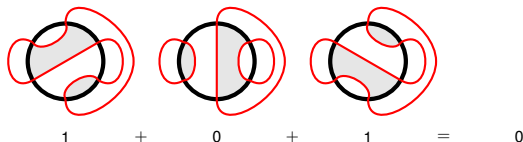
If $\langle \cdot | \cdot \rangle$ is to be nondegenerate, any three chord diagrams related by bypass surgery should sum to 0: *bypass relation*.



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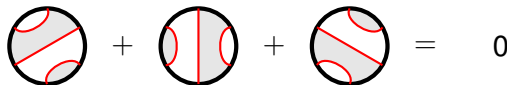
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So we define a vector space

$$V_n = \frac{\mathbb{Z}_2 \langle \text{Chord diagrams with } n \text{ chords} \rangle}{\text{Bypass relation}}$$

A vector space of chord diagrams

Theorem (M.)

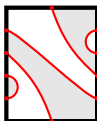
V_n has dimension 2^{n-1} and the diagrams from chessboards of $n - 1$ squares form a basis.

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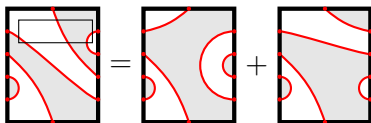


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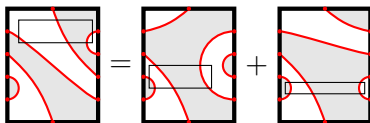


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E.g.

$$\begin{aligned}
 & \text{Chessboard with } 2 \times 1 \text{ shaded squares} = \text{Chessboard with } 1 \times 1 \text{ shaded squares} + \text{Chessboard with } 1 \times 1 \text{ shaded squares} \\
 & = \text{Chessboard with } 1 \times 1 \text{ shaded squares} + \text{Chessboard with } 1 \times 1 \text{ shaded squares} + \text{Chessboard with } 1 \times 1 \text{ shaded squares} + \text{Chessboard with } 1 \times 1 \text{ shaded squares} \\
 & = \Gamma_{ppqq} + \Gamma_{pqqp} + \Gamma_{appq} + \Gamma_{qpqp}
 \end{aligned}$$

Outline

- 1 Introduction
- 2 Combinatorial and algebraic structure
- 3 Contact topology
 - Chord diagrams and contact structures
 - Bypasses
 - Contact QFT = Quantum pawn dynamics
- 4 Holomorphic invariants

Chord diagrams and contact structures

Giroux (1991): theory of *convex surfaces*.

A chord diagram Γ / *dividing set* on a disc D describes a contact structure ξ_Γ on a neighbourhood $D \times I$ of D .

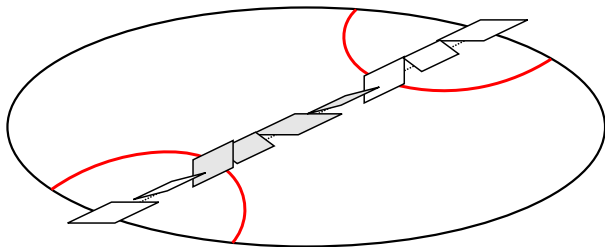
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Roughly speaking, the contact planes are

- Tangent to ∂D
- “Perpendicular” to D precisely along Γ



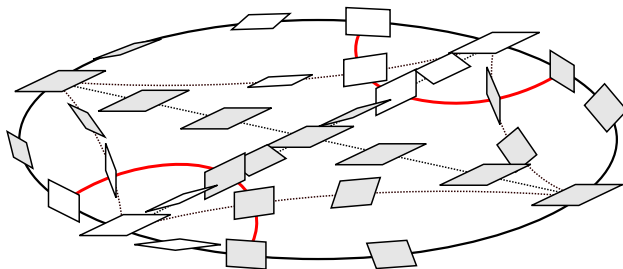
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Colours in chord diagram = visible side of contact plane.

Overtwisted contact structures

Eliashberg (1989): fundamentally 2 types of contact structures.

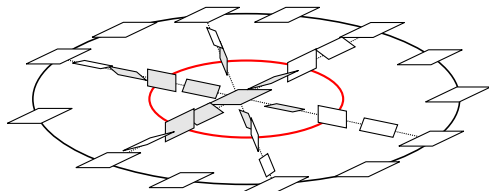
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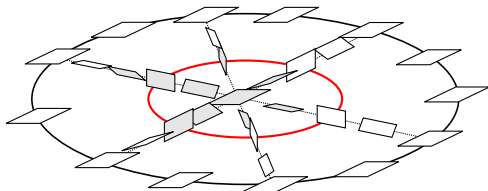


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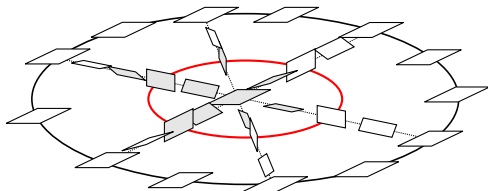
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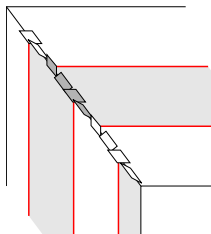
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- Overtwisted contact geometry reduces to (well-understood) homotopy theory. Tight contact structures offer important topological information.
- Eliashberg (1992): contact structure near an S^2 is tight iff dividing set is *connected*. If so, contact structure extends uniquely (up to isotopy) to a tight contact structure on B^3 .

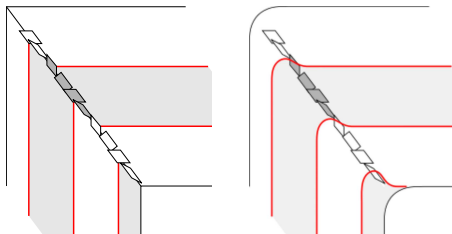
Contact corners

When two convex surfaces meet along a boundary, contact planes are arranged as shown.



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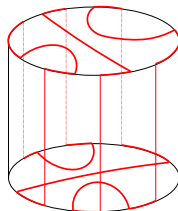
Proposition

Let Γ_0, Γ_1 be chord diagrams. The following are equivalent:

- $\langle \Gamma_0 | \Gamma_1 \rangle = 1$.
- The solid cylinder with dividing set Γ_0 on the bottom and Γ_1 on the top has a tight contact structure.

Bypasses

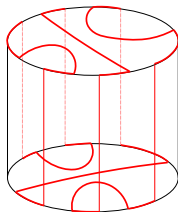
Honda (2000's): any 3-manifold can be built up from a surface and dividing set by adding *bypasses*.



Effect on dividing set is “bypass surgery” as defined earlier.

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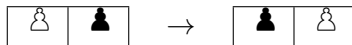
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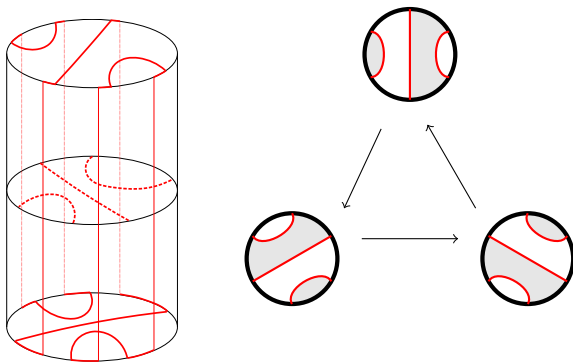
$$\langle \Gamma_{pq} | \Gamma_{qp} \rangle = 1$$

or



Bypasses

Stacking two bypasses on top of each other produces an overtwisted contact structure!



Can build something like a *triangulated category* out of dividing sets and contact structures (Honda, M.). V_n is the *Grothendieck group*.

Contact TQFT = Quantum pawn dynamics

These definitions give many of the properties of a (2+1)-dimensional *topological quantum field theory*.

- Contact structure near disc (2-dim) \rightsquigarrow “states” in V_n
- Contact structure over cylinder (2+1-dim) \rightsquigarrow element of \mathbb{Z}_2 .
- “Possibility of a tight contact structure from one state to another” \rightsquigarrow inner product $\langle \cdot | \cdot \rangle : V_n \otimes V_n \longrightarrow \mathbb{Z}_2$.

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Theorem (M.)

“Contact TQFT is isomorphic to quantum pawn dynamics.”

Outline

- 1 Introduction
- 2 Combinatorial and algebraic structure
- 3 Contact topology
- 4 Holomorphic invariants
 - Sutured Floer homology
 - A “computation”

Sutured Floer homology

Actually all the above comes from *sutured Floer homology*, a holomorphic invariant of sutured manifolds.

Very roughly... (Ozsváth–Szabó 2004, Juhasz 2006)

- A *sutured manifold* is a 3-manifold M with boundary, and some curves Γ on ∂M dividing ∂M into alternating positive and negative regions.

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- Consider $\Sigma \times I \times \mathbb{R}$ as a symplectic manifold with an almost complex structure and consider holomorphic curves

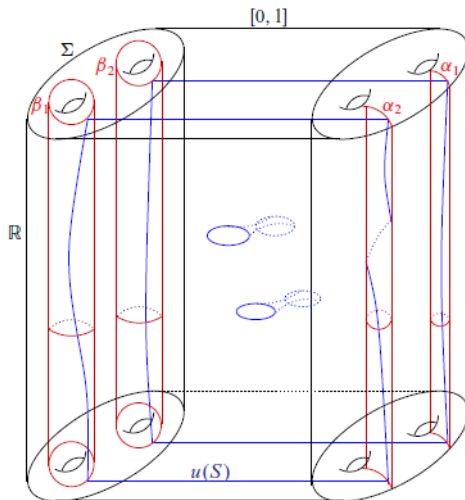
$$u : S \longrightarrow \Sigma \times I \times \mathbb{R}$$

where S is a Riemann surface.

- Boundary conditions based on Heegaard curves α_i and β_i .

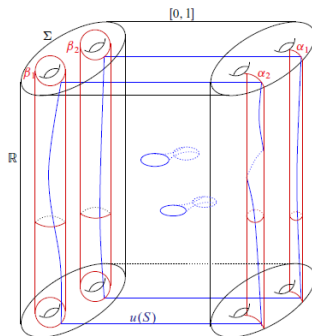
Sutured Floer homology

Cylindrical picture of Lipshitz (2006):



Sutured Floer homology

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$$\text{ind}(D\bar{\partial}) = k - \chi(S) + \sum_{i=1}^k \mu(a_i) - \sum_{i=1}^k \mu(b_i).$$

Sutured Floer homology

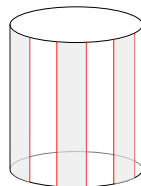
- Chain complex generated by boundary conditions, which are *intersections* of boundary curves.

$$z_1 \in \alpha_1 \cap \beta_{\sigma(1)}, z_2 \in \alpha_2 \cap \beta_{\sigma(2)}, \dots, z_k \in \alpha_k \cap \beta_{\sigma(k)}.$$

- Differential counting index-1 holomorphic curves between boundary conditions.
- Resulting homology is $SFH(M, \Gamma)$.
- Etnyre–Honda (2009): Any *contact structure* ξ on (M, Γ) defines a natural *element* $c(\xi) \in SFH(M, \Gamma)$.

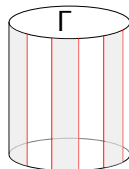
Solid tori

We consider the *sutured solid torus*
 $D^2 \times S^1$ with $2n$ longitudinal curves
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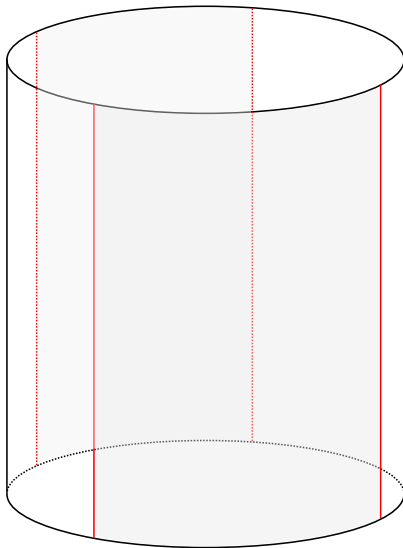


Theorem (M.)

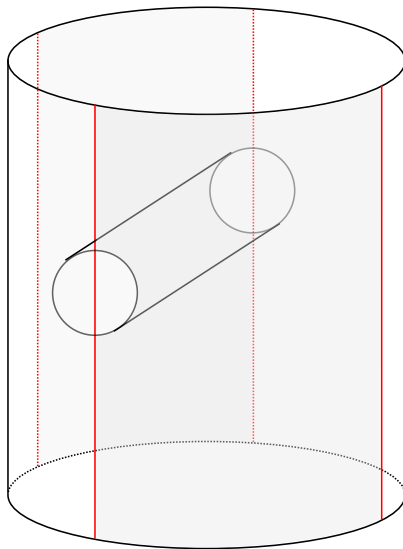
$$SFH(D^2 \times S^1, F_n \times S^1) \cong V_n = \frac{\mathbb{Z}_2 \langle \text{Chord diagrams w/ } n \text{ chords} \rangle}{\text{Bypass relation}}$$

Any chord diagram Γ in V_n corresponds to a contact structure ξ_Γ on $D^2 \times S^1$ and maps to $c(\xi_\Gamma)$.

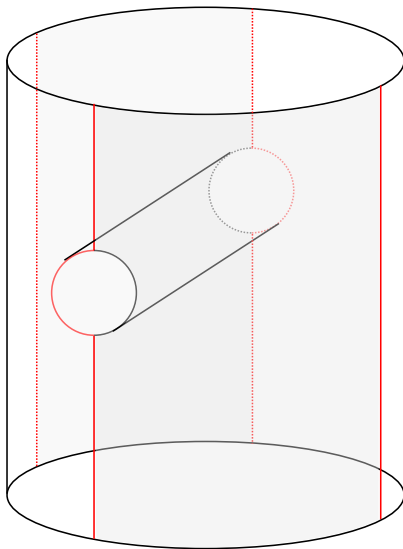
A “computation” of Sutured Floer homology



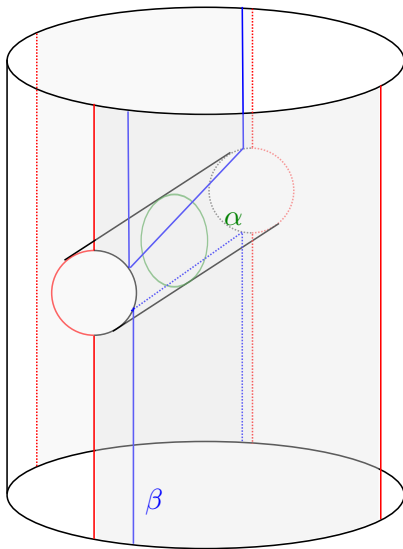
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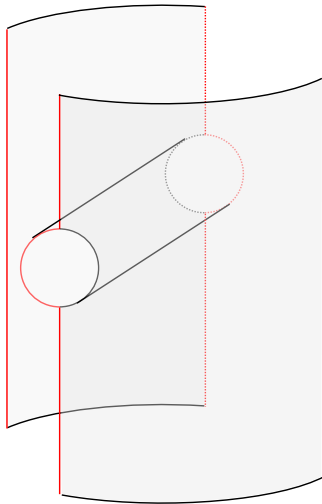
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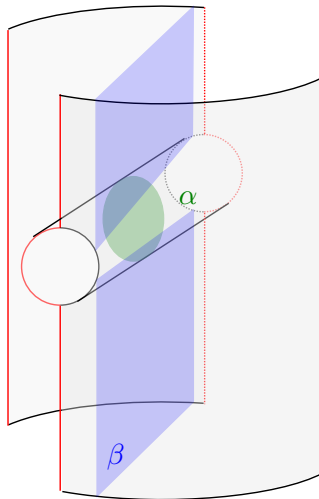
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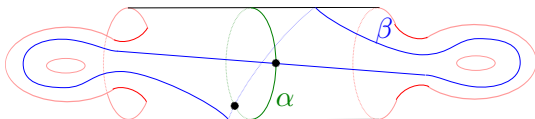
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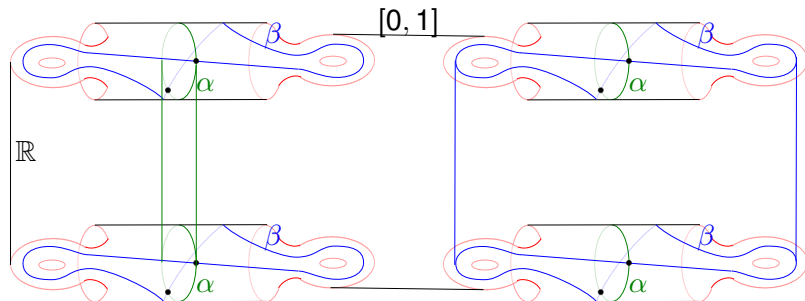


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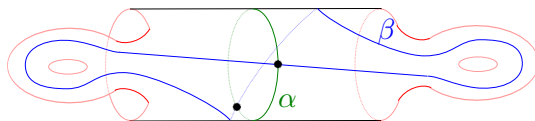
Chain complex = $\mathbb{Z}_2 \oplus \mathbb{Z}_2$.

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Chain complex = $\mathbb{Z}_2 \oplus \mathbb{Z}_2$.

Nowhere for holomorphic curves to go! $\partial = 0$.

$$SFH = \mathbb{Z}_2 \oplus \mathbb{Z}_2 = V_2$$

Thanks for listening!

References:

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- D. Mathews, *Sutured TQFT, torsion, and tori* (2011) arXiv 1102.3450.
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- D. Mathews, *Contact topology and holomorphic invariants via elementary combinatorics* (2012) arXiv 1212.1759.