

Contact topology and holomorphic invariants via elementary combinatorics

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Outline

- 1 Introduction
- 2 Background
- 3 Quantum Pawn Dynamics (QPD)
- 4 Chord diagrams
- 5 String homology
- 6 Holomorphic invariants

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Overview

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- Much of it is quite involved, but *in the simplest cases* some of this structure reduces to some *elementary combinatorics and algebra* which is interesting in its own right.

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- Give some *very* brief background on symplectic / contact geometry and holomorphic curves.
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 - “Quantum pawn dynamics”
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 - “String homology”
- Explain how this elementary combinatorics arises from holomorphic invariants.

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 - Symplectic geometry
 - Contact geometry
 - Complex structures
 - Holomorphic curves
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Symplectic manifolds

Definition

A symplectic manifold is a pair

$$(M, \omega)$$

where

- *M is a smooth manifold*
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Such M must be even-dimensional.

Symplectic manifolds

Main example:

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Structure of Hamiltonian classical mechanics:

- Given a smooth function $H : M \rightarrow \mathbb{R}$ (Hamiltonian) we obtain a 1-form dH and a dual vector field X_H via

$$\omega(X_H, \cdot) = dH$$

Contact geometry

“The odd-dimensional sibling of symplectic geometry”

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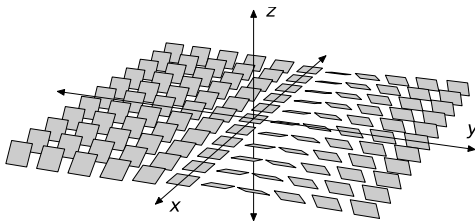
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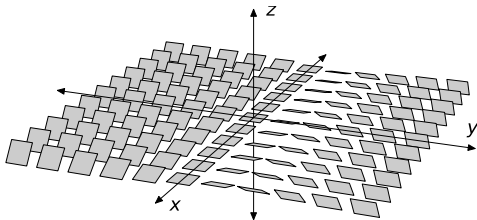


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A contact structure is locally the kernel of a *contact 1-form* α .

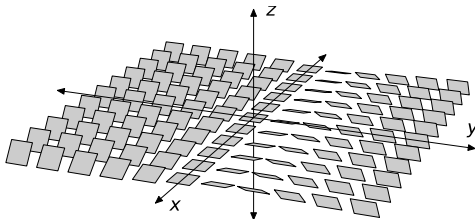
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- Non-integrability means $\alpha \wedge (d\alpha)^n \neq 0$ everywhere.
- (So $\alpha|_{\xi}$ is a symplectic form at each point.)

Contact geometry

Main example of contact manifold:

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Today: $n = 1$, 3-dimensional contact geometry.

- Much of 3-dimensional contact geometry can be described **combinatorially**.

Symplectic vs complex geometry

- Complex geometry also only exists in *even* number of dimensions.
- Gromov (1985): Consider *almost complex structures* on symplectic manifolds.

Symplectic vs complex geometry

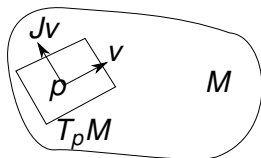
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Definition

An *almost complex structure* on a smooth manifold is a map

$$J : TM \longrightarrow TM$$

preserving each fibre $T_p M$ and satisfying $J^2 = -1$.



Almost complex vs complex

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Existence:

- Not every symplectic manifold has a complex structure...
- ... but every symplectic manifold has a compatible almost complex structure J , and all choices of J are homotopic.

(Compatible: J and ω behave in linear algebra like i and $dx \wedge dy$. $\omega(v, w) = \omega(Jv, Jw)$ and $\omega(v, Jv) > 0$)

Holomorphic curves

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Holomorphic curves

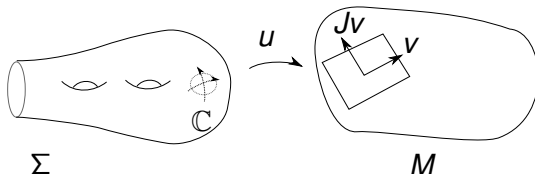
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A *holomorphic curve* is a map $u : \Sigma \rightarrow M$, where Σ is a Riemann surface, satisfying the Cauchy-Riemann equations

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Holomorphic curves

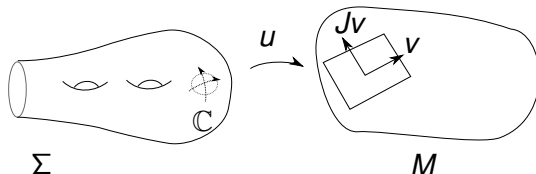
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An *almost* complex structure is sufficient for the C-R equations.

Moduli spaces

- Given appropriate constraints (marked points, boundary conditions) and transversality, the space of holomorphic curves is a finite-dimensional orbifold: *moduli space* \mathcal{M} .

Moduli spaces

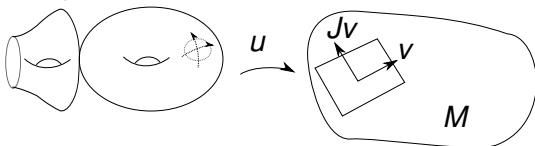
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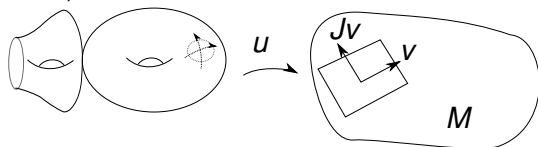
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- \mathcal{M} and $\overline{\mathcal{M}}$ encode a great deal of information about M .
- Some powerful invariants use only the *codimension-1 boundary* of $\overline{\mathcal{M}}$.

Floer Homology

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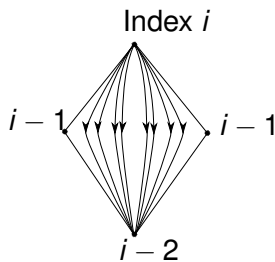
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Analogy: singular homology via Morse complex.

- Complex generated by critical points of Morse function f .
- ∂ counts 0-dimensional families of trajectories of ∇f .



The power of holomorphic invariants

Floer homology theories give very powerful invariants of 3-manifolds, knots, etc...

- Related to Seiberg–Witten theory, Donaldson–Thomas theory, etc...
- E.g., *knot Floer homology* can compute the genus of a knot.

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- Related to Seiberg–Witten theory, Donaldson–Thomas theory, etc...
- E.g., *knot Floer homology* can compute the genus of a knot.
- For a *less complicated* variant called *sutured Floer homology*, and a *simple class* of manifolds $M = \Sigma \times S^1$, we obtain all the combinatorial structure we are about to see, and more...

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 - Pawns and anti-pawns
 - Creation and annihilation operators
 - Adjoints
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Quantum Pawn Dynamics

- Pawns on a finite 1-dimensional chessboard.
- A state of the QPD universe:

$$w = \begin{array}{|c|c|c|c|c|c|} \hline \text{♙} & & \text{♙} & \text{♙} & & \\ \hline \end{array}$$

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Quantum pawns: “Inner product” $\langle \cdot | \cdot \rangle$ describes the possibility of pawn moves from one state to another.

- Valued in \mathbb{Z}_2 .

Definition (Pawn “inner product”)

$$\langle w_0 | w_1 \rangle = \begin{cases} 1 & \text{if it is possible for pawns to move from } w_0 \text{ to } w_1 \\ & \text{(this includes the case } w_0 = w_1); \\ 0 & \text{if not.} \end{cases}$$

Quantum Pawn Dynamics

E.g.

$$\langle \begin{array}{|c|c|c|c|c|c|} \hline \text{pawn} & & \text{pawn} & \text{pawn} & & \\ \hline \end{array} \mid \begin{array}{|c|c|c|c|c|c|} \hline & \text{pawn} & \text{pawn} & & \text{pawn} & \\ \hline \end{array} \rangle = 1$$

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Also, entangled chessboards.

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Note asymmetry of $\langle \cdot \mid \cdot \rangle$.

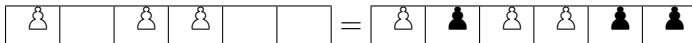
A “booleanized” partial order. (Complete lattice.)

Dirac Pawn Sea

- Introduce the *anti-pawn* = *absence of pawn*.

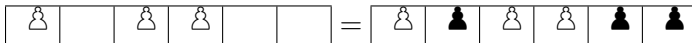
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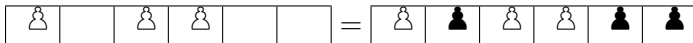
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- A pawn moving right is equivalent to an anti-pawn moving left.
- Let
 - n_p = number of pawns
 - n_q = numbers of anti-pawns
 - n = number of squares on board = $n_p + n_q$

Creation and annihilation operators

The *initial pawn creation operator* $a_{p,0}^*$ adjoins a new *initial* (leftmost) square to the chessboard, containing a pawn.

$$a_{p,0}^* \begin{array}{|c|c|c|c|c|c|} \hline \text{♔} & \text{♚} & \text{♔} & \text{♔} & \text{♚} & \text{♚} \\ \hline \end{array} = \begin{array}{|c|c|c|c|c|c|c|} \hline \text{♔} & \text{♔} & \text{♚} & \text{♔} & \text{♔} & \text{♚} & \text{♚} \\ \hline \end{array}$$

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The *initial pawn annihilation operator* $a_{p,0}$ deletes the leftmost square from the chessboard, and a pawn on it.

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If no pawn (anti-pawn) in the leftmost square, try to delete...

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If no pawn (anti-pawn) in the leftmost square, try to delete...

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Creation and annihilation operators

The *initial pawn creation operator* $a_{p,0}^*$ adjoins a new *initial* (leftmost) square to the chessboard, containing a pawn.

$$a_{p,0}^* \begin{array}{|c|c|c|c|c|c|} \hline \text{white} & \text{black} & \text{white} & \text{white} & \text{black} & \text{black} \\ \hline \end{array} = \begin{array}{|c|c|c|c|c|c|c|} \hline \text{white} & \text{white} & \text{black} & \text{white} & \text{white} & \text{black} & \text{black} \\ \hline \end{array}$$

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Similar initial anti-pawn annihilation $a_{q,0}$ and creation $a_{q,0}^\dagger$.

Creation of chessboards

- The *vacuum* state of the QPD universe is the null chessboard \emptyset .
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- The $*$ and † refer to *adjoints* — just as in quantum field theory.
(Actually they form a *Galois connection* on partial orders.)

Adjoint

- Recall an adjoint f^* of an operator f usually means that

$$\langle fx|y\rangle = \langle x|f^*y\rangle, \quad \langle x|fy\rangle = \langle f^*x|y\rangle.$$

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- As our “inner product” is asymmetric, we have *two distinct adjoints* f^*, f^\dagger of an operator f .

$$\langle fx|y\rangle = \langle x|f^*y\rangle, \quad \langle x|fy\rangle = \langle f^\dagger x|y\rangle.$$

So $f^{*\dagger} = f^{\dagger*} = f$.

Initial creation and annihilation are adjoint

Proposition

$$\langle a_{p,0}x|y\rangle = \langle x|a_{p,0}^*y\rangle$$

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Similarly, initial anti-pawn creation/annihilation \dagger -adjoint.

Adjoining adjoints

What is $a_{p,0}^{**}$? What operator f satisfies

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Call this operator $a_{p,1}$.

Adjoining adjoints

Keep going. What is $a_{p,0}^{***} = a_{p,1}^*$? What operator g satisfies

$$\langle a_{p,1}x|y\rangle = \langle x|gy\rangle?$$

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(I.e. inserts an extra square with pawn next to it)

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 \end{aligned}$$

Iterated adjoints

Proposition

The iterated adjoints of $a_{p,0}$ are

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Similarly for anti-pawns in the opposite direction.

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(A simplicial structure.)

Adjoint periodicity

Hence

$$a_{p,0}^{*2n_p+2} = a_{p,\Omega}$$

where n_p = number of pawns.

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One can also show that the *duality* operator defined by

$$\langle u|v \rangle = \langle v|Hu \rangle$$

satisfies

Theorem (M.)

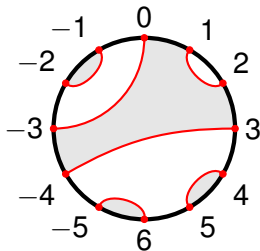
$$H^{2n+2} = 1.$$

Outline

- 1 Introduction
- 2 Background
- 3 Quantum Pawn Dynamics (QPD)
- 4 Chord diagrams
 - Chord diagrams
 - Creation and annihilation
 - Chessboards and chord diagrams
 - “Inner product” on chord diagrams
 - Bypasses
 - Chord diagram vector space
- 5 String homology

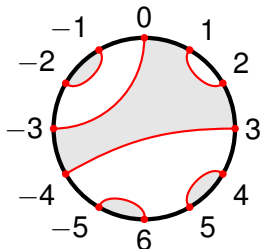
Chord diagrams

Consider a disc D with some $2n + 2$ points F marked on ∂D .
 A *chord diagram* is a collection of non-intersecting curves on D joining points of F .
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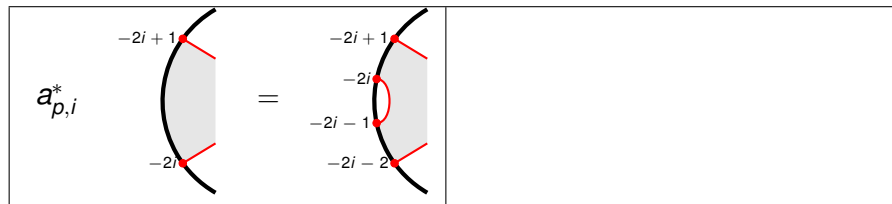
- Curves join points of opposite parity, so shade as shown.
- 0 is a basepoint.
- Label points mod $2n + 2$.

Creation and annihilation of chords

Define *creation operators* $a_{p,i}^*$, $a_{q,i}^\dagger$ to insert a new chord in a specific place in a chord diagram as shown.

Creation and annihilation of chords

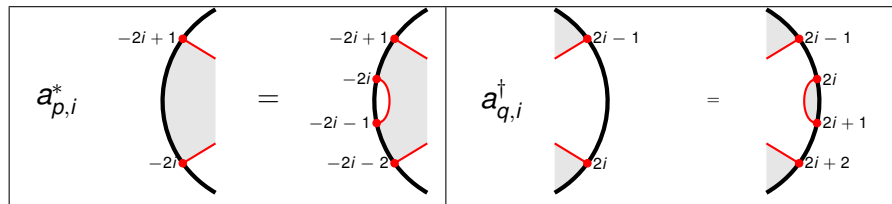
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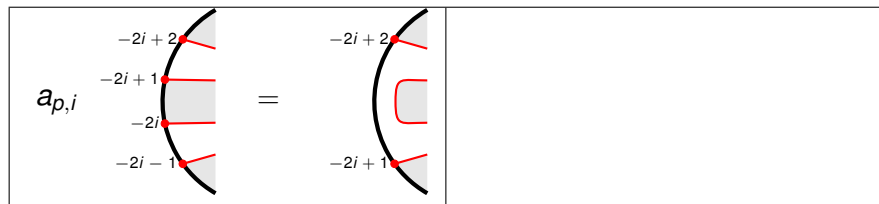
$a_{q,i}^\dagger$ creates a *black* region i spots down on the right.

Creation and annihilation of chords

Define *annihilation operators* $a_{p,i}$, $a_{q,i}$ to **close off chords** in a chord diagram as shown.

Creation and annihilation of chords

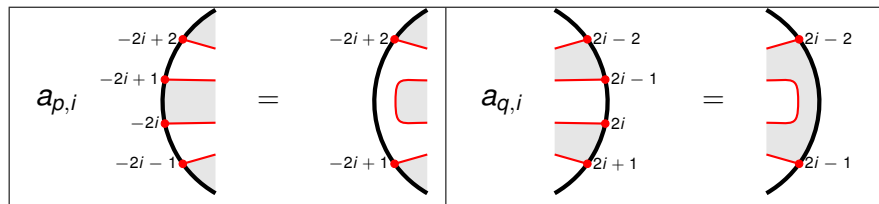
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Diagrams of chessboards

The simplest chord diagram is called the *vacuum* Γ_\emptyset .



Build up more complicated diagrams with creation operators.

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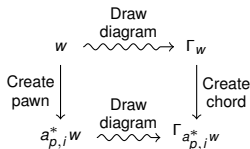
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Build up more complicated diagrams with creation operators.

Proposition (M.)

*For any chessboard w , there is a chord diagram Γ_w (called a **slalom** chord diagram) such that creation and annihilation operators agree (are equivariant): $\Gamma_{a_{p,i}^* w} = a_{p,i}^* \Gamma_w$.*



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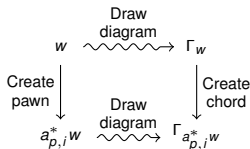
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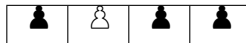
- If chessboard has n squares, i.e. $|w| = n$, then Γ_w has $n + 1$ chords.

Ski slopes

Construction of the *slalom skiing* chord diagram of a chessboard.

qpqq

\leftrightarrow

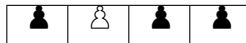


Ski slopes

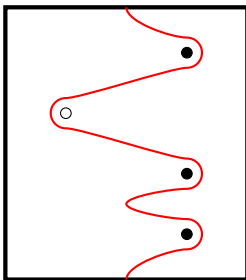
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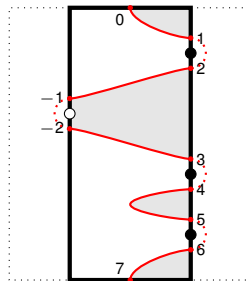
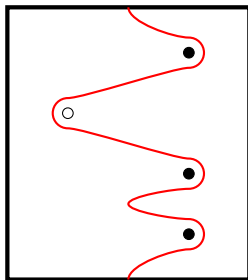
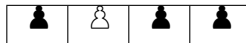


\leftrightarrow



Ski slopes

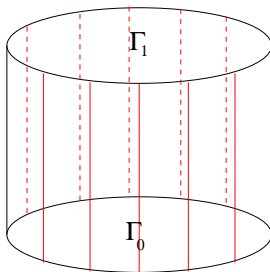
Construction of the *slalom skiing* chord diagram of a chessboard.



$$= \Gamma_{qpqq}$$

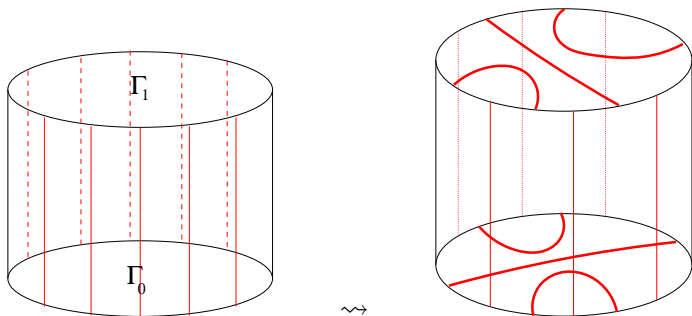
An “Inner product” on chord diagrams

There's a bilinear form on chord diagrams defined by *entering into a cylinder*.



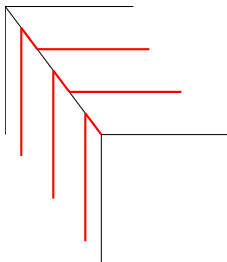
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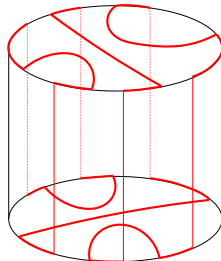
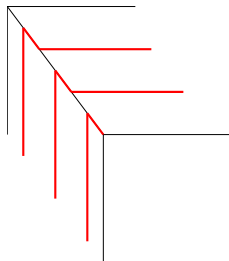
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Note curves don't meet at corners! We treat corners as shown.



An “Inner product” on chord diagrams

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Definition

$$\langle \Gamma_0 | \Gamma_1 \rangle = \begin{cases} 1 & \text{if the resulting curves on the cylinder} \\ & \text{form a single connected curve;} \\ 0 & \text{if the result is disconnected.} \end{cases}$$

Theorem (M.)

For any two chessboards w_0, w_1 ,

$$\langle w_0 | w_1 \rangle = \langle \Gamma_{w_0} | \Gamma_{w_1} \rangle.$$

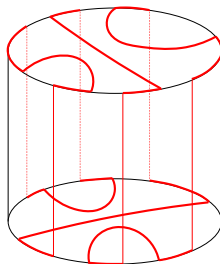
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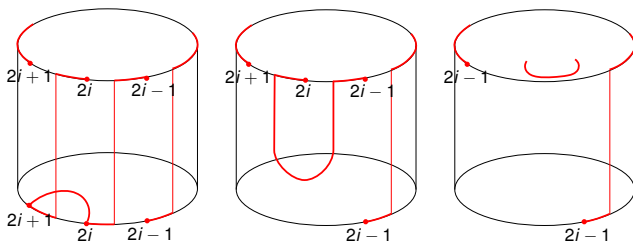
E.g.

$$\left\langle \begin{array}{|c|c|} \hline \text{white} & \text{black} \\ \hline \end{array} \mid \begin{array}{|c|c|} \hline \text{black} & \text{white} \\ \hline \end{array} \right\rangle = 1 =$$



Adjoint

Adjoint relations can be seen *topologically* as “finger moves”.



$$\langle a_{q,i}^\dagger \Gamma_0 \mid \Gamma_1 \rangle = \langle \Gamma_0 \mid a_{q,i} \Gamma_1 \rangle$$

Now perhaps believable that adjoint is periodic.

Bypass surgery

In a chord diagram on disc D , consider a sub-disc B as shown:



Bypass surgery

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Two natural ways to adjust this chord diagram, consistent with the colours: *bypass surgeries*.


 Γ'

 Γ

 Γ''

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Proposition

With $\Gamma, \Gamma', \Gamma''$ as above, for any Γ_1 ,

$$\langle \Gamma | \Gamma_1 \rangle + \langle \Gamma' | \Gamma_1 \rangle + \langle \Gamma'' | \Gamma_1 \rangle = 0.$$

Bypass surgery

Idea of proof:

The diagram shows three identical circular disks arranged horizontally. Each disk has a black diagonal chord from the top-left to the bottom-right. The region to the left of the chord is shaded gray. Red loops are drawn on each disk, representing a specific chord diagram. Below the disks, the equation $1 + 0 + 1 = 0$ is written, where the numbers 1, 0, and 1 correspond to the three disks respectively.

$$1 + 0 + 1 = 0$$

Bypass surgery

Idea of proof:

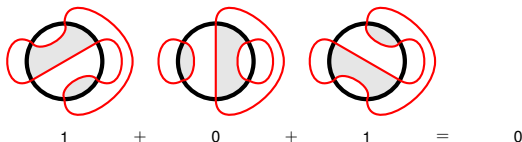
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If $\langle \cdot | \cdot \rangle$ is to be nondegenerate, any three chord diagrams related by bypass surgery should sum to 0: *bypass relation*.

$$+ + = 0$$

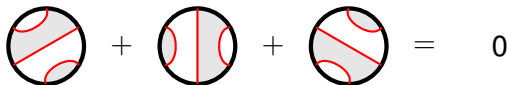
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$$+ + = 0$$

So we define a vector space

$$V_n = \frac{\mathbb{Z}_2 \langle \text{Chord diagrams with } n+1 \text{ chords} \rangle}{\text{Bypass relation}}$$

A vector space of chord diagrams

Theorem (M.)

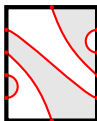
V_n has dimension 2^n and the slalom diagrams from chessboards of n squares form a basis.

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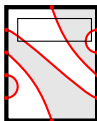


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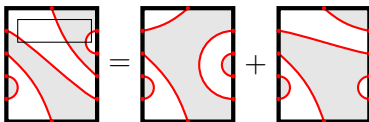


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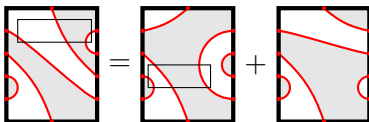


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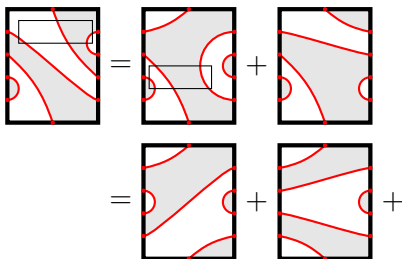


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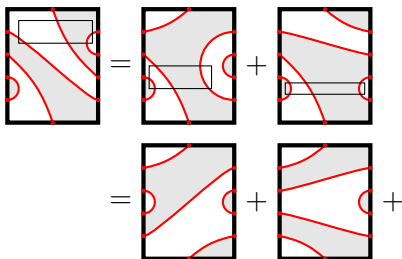


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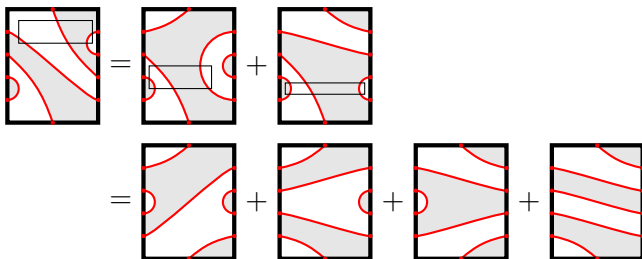


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E.g.

The diagram illustrates the decomposition of a chessboard with a single square (represented by a white square in the top-left corner of a larger square) into four slalom diagrams. The first row shows the decomposition: the chessboard equals the sum of two diagrams. The second row shows these two diagrams further decomposed into four diagrams. The third row shows the corresponding algebraic expression using the Γ notation.

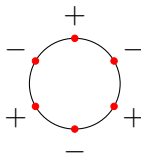
$$\begin{aligned}
 & \text{Chessboard} = \text{Diagram 1} + \text{Diagram 2} \\
 & \text{Diagram 1} = \text{Diagram 3} + \text{Diagram 4} \\
 & \text{Diagram 2} = \text{Diagram 5} + \text{Diagram 6} \\
 & = \Gamma_{ppqq} + \Gamma_{pqqp} + \Gamma_{appq} + \Gamma_{qpqp}
 \end{aligned}$$

Outline

- 1 Introduction
- 2 Background
- 3 Quantum Pawn Dynamics (QPD)
- 4 Chord diagrams
- 5 String homology
 - The string complex
 - Calculation of homology
- 6 Holomorphic invariants

String diagrams

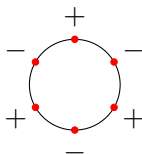
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The points of F are **signed**: half $+$, half $-$.



We usually consider points F which **alternate in sign**.
(Points were effectively oriented previously...)

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Definition

A **string diagram** s is a collection of oriented immersed curves on D with $\partial s = F$.

The string complex

We define a **chain complex** based on string diagrams.

Definition

$$\widehat{CS}(D, F) = \frac{\mathbb{Z}_2 \langle \text{String diagrams on } (D, F) \rangle}{\mathbb{Z}_2 \langle \text{String diagrams with contractible curves} \rangle}$$

I.e. formal sums of string diagrams; contractible curves = 0.

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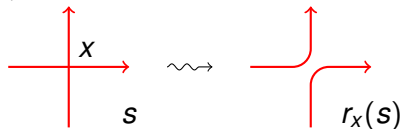
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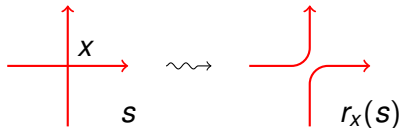
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Definition

$$\partial s = \sum_{x \text{ crossing of } s} r_x(s).$$

Well-definition

Is ∂ well defined?

Show that ∂ is unchanged by “string Reidemeister moves”. E.g.

Well-definition

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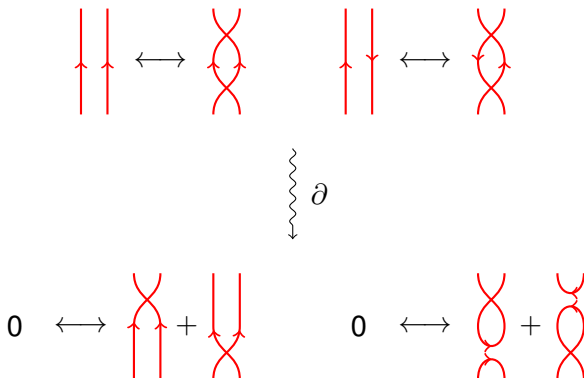
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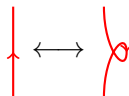
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This shows why mod 2 is useful...

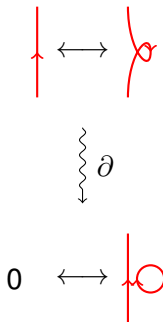
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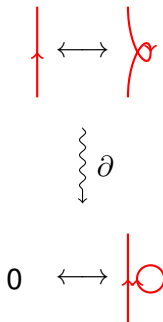
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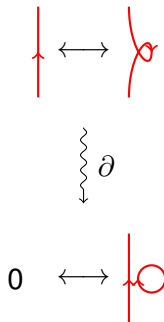
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What is the homology $\widehat{HS}(D, F)$?

Calculation of homology

Note $\widehat{CS}(D, F)$ is well-defined, whether F is alternating or not.

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For alternating F , $\widehat{HS}(D, F)$ is generated by chord diagrams, and the bypass relation is satisfied. In fact,

$$\widehat{HS}(D^2, F) \cong \frac{\mathbb{Z}_2 \langle \text{Chord diagrams on } (D, F) \rangle}{\text{Bypass relation}} \cong V_n.$$

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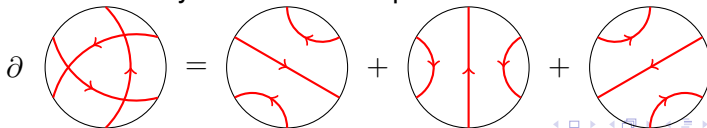
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A “reason” why the theorem is plausible:



Proof of proposition

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Consider the **switching** operation on string diagrams

$$W : \widehat{CS}(\Sigma, F) \longrightarrow \widehat{CS}(\Sigma, F).$$



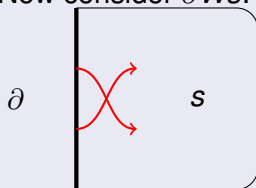
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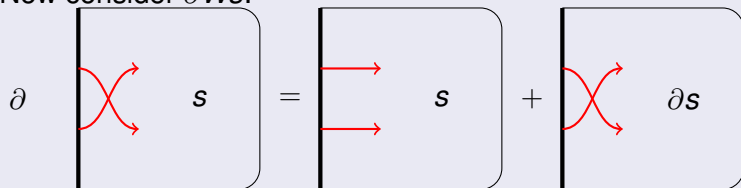
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Proof of proposition

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Diagram illustrating the boundary operation ∂ on the switching operation W applied to a string diagram s :

Left side: A string diagram with two horizontal lines labeled q and p from top to bottom. A vertical line is to the left of the lines. The lines q and p cross each other. The entire diagram is enclosed in a rounded rectangle labeled s . A vertical line is to the left of the rectangle. A wavy arrow labeled ∂ points to the right.

Right side: A string diagram with two horizontal lines labeled q and p from top to bottom. A vertical line is to the left of the lines. The lines q and p are parallel. The entire diagram is enclosed in a rounded rectangle labeled s . A vertical line is to the left of the rectangle. A wavy arrow labeled ∂ points to the right.

Equation:

$$\partial Ws = s + W\partial s$$

Proof of proposition

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Now consider ∂Ws :

$$\partial Ws = s + W\partial s$$

Thus W is a chain homotopy from 1 to 0, and $\widehat{HS}(\Sigma, F) = 0$.

Outline

- 1 Introduction
- 2 Background
- 3 Quantum Pawn Dynamics (QPD)
- 4 Chord diagrams
- 5 String homology
- 6 Holomorphic invariants
 - Sutured Floer homology
 - A “computation”
 - Contact invariants
 - Bypasses

Sutured Floer homology

Actually all the above comes from *sutured Floer homology*, a holomorphic invariant of sutured manifolds.

Very roughly... (Ozsváth–Szabó 2004, Juhasz 2006)

- A *sutured manifold* is a 3-manifold M with boundary, and some curves Γ on ∂M dividing ∂M into alternating positive and negative regions.

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- Given (M, Γ) , take a *Heegaard decomposition* with surface Σ and curves $\alpha_1, \dots, \alpha_k$ bounding discs on one side and β_1, \dots, β_k bounding discs on the other.

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- Consider $\Sigma \times I \times \mathbb{R}$ as a symplectic manifold with an almost complex structure and consider holomorphic curves

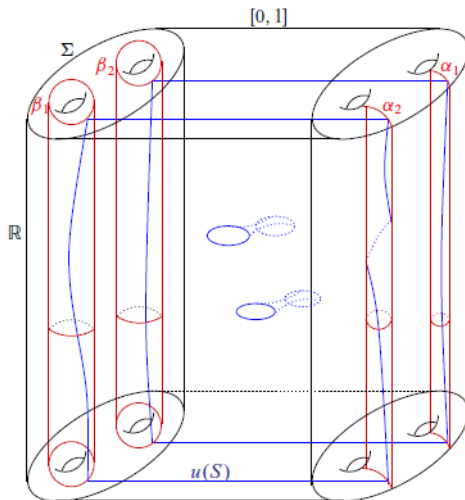
$$u : S \longrightarrow \Sigma \times I \times \mathbb{R}$$

where S is a Riemann surface.

- Boundary conditions based on Heegaard curves α_i and β_i .

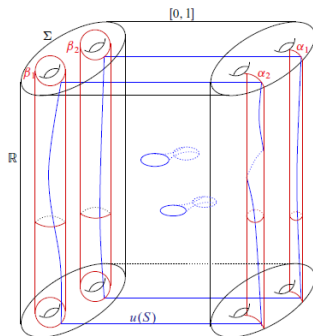
Sutured Floer homology

Cylindrical picture of Lipshitz (2006):



Sutured Floer homology

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$$\text{ind}(D\bar{\partial}) = k - \chi(S) + \sum_{i=1}^k \mu(a_i) - \sum_{i=1}^k \mu(b_i).$$

Sutured Floer homology

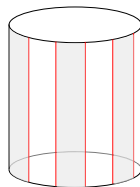
- Chain complex generated by boundary conditions, which are *intersections* of boundary curves.

$$z_1 \in \alpha_1 \cap \beta_{\sigma(1)}, z_2 \in \alpha_2 \cap \beta_{\sigma(2)}, \dots, z_k \in \alpha_k \cap \beta_{\sigma(k)}.$$

- Differential counting index-1 holomorphic curves between boundary conditions.
- Resulting homology is $SFH(M, \Gamma)$.
- Etnyre–Honda (2009): Any *contact structure* ξ on (M, Γ) defines a natural *element* $c(\xi) \in SFH(M, \Gamma)$.

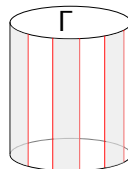
Solid tori

We consider the *sutured solid torus* $D^2 \times S^1$ with $2n + 2$ longitudinal curves $F \times S^1$. ($|F| = 2n + 2$)



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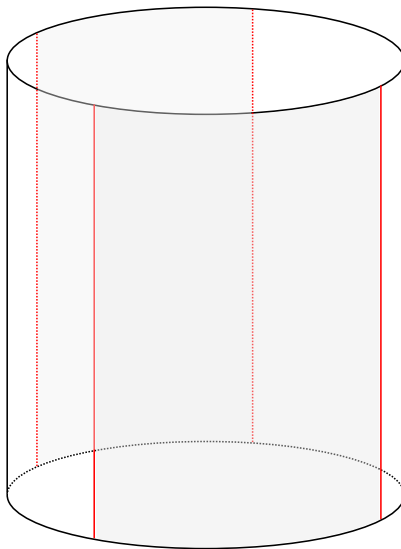
Theorem (M.)

$$SFH(D^2 \times S^1, F \times S^1) \cong V_n = \frac{\mathbb{Z}_2 \langle \text{Chord diagrams w/ } n+1 \text{ chords} \rangle}{\text{Bypass relation}}$$

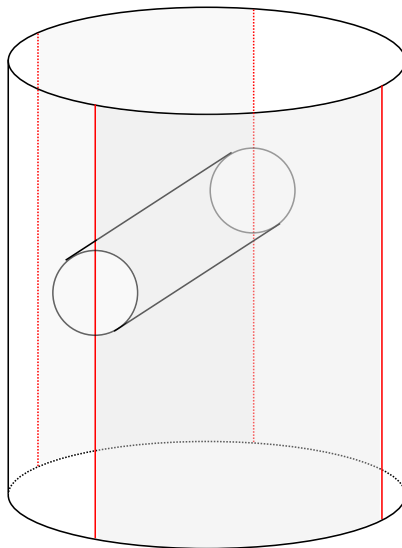
$$\cong \widehat{HS}(D^2, F)$$

Any chord diagram Γ in V_n or $\widehat{HS}(D^2, F)$ corresponds to a contact structure ξ_Γ on $D^2 \times S^1$ and maps to $c(\xi_\Gamma)$.

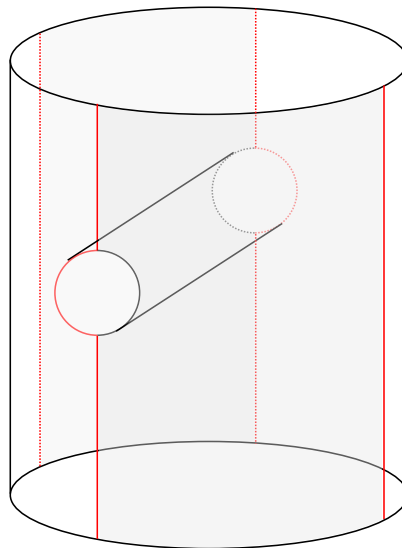
A “computation” of Sutured Floer homology



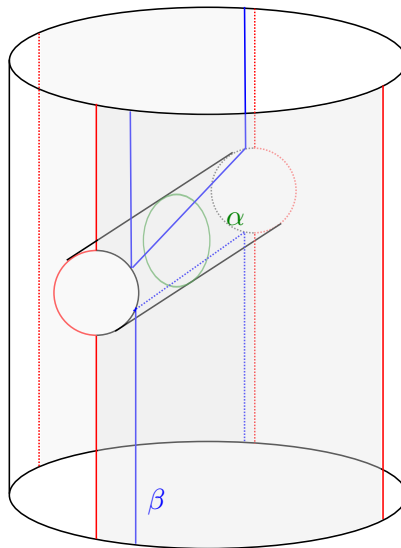
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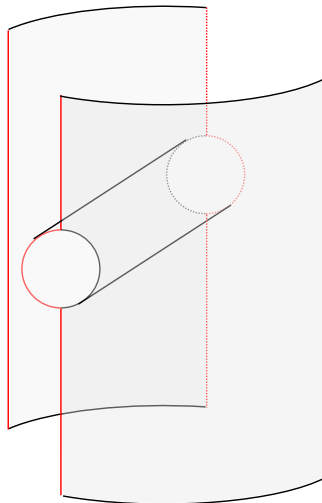
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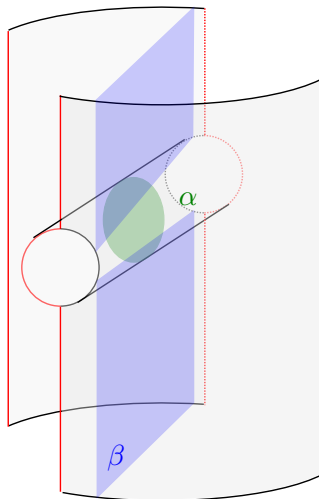
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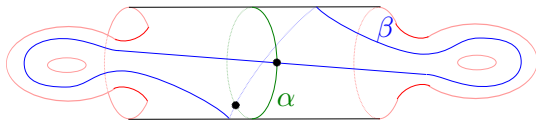
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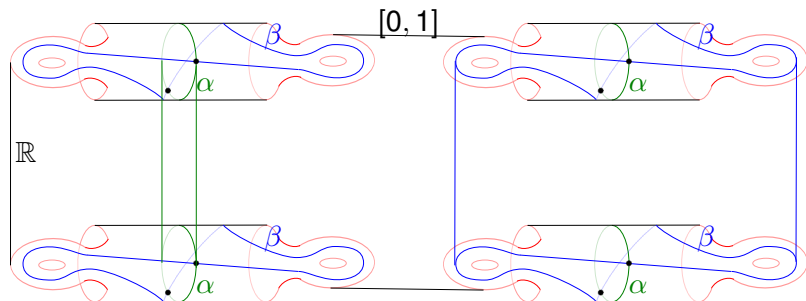


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Chain complex = $\mathbb{Z}_2 \oplus \mathbb{Z}_2$.

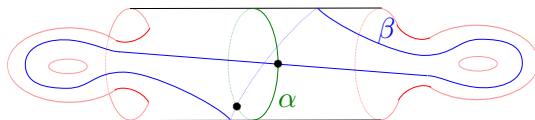
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Where do holomorphic curves go?

A “computation” of Sutured Floer homology



Chain complex = $\mathbb{Z}_2 \oplus \mathbb{Z}_2$.

Where do holomorphic curves go?

Nowhere for holomorphic curves to go! $\partial = 0$.

$$SFH = \mathbb{Z}_2 \oplus \mathbb{Z}_2 = V_1$$

Chord diagrams and contact structures

Giroux (1991): theory of *convex surfaces*.

A chord diagram Γ / *dividing set* on a disc D describes a contact structure ξ_Γ on a neighbourhood $D \times I$ of D .

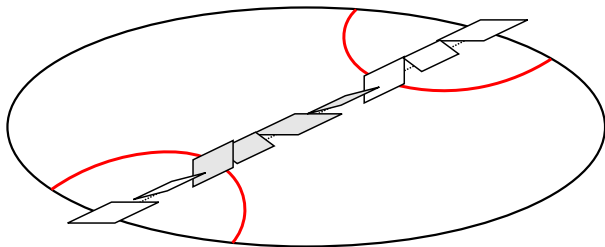
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Roughly speaking, the contact planes are

- Tangent to ∂D
- “Perpendicular” to D precisely along Γ



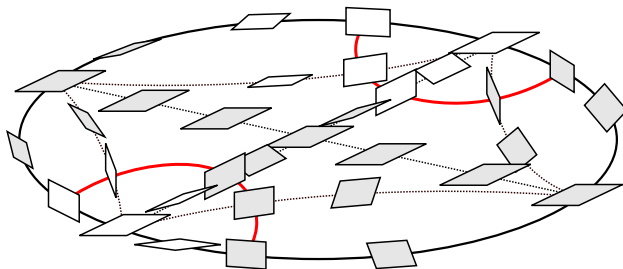
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Colours in chord diagram = visible side of contact plane.

Overtwisted contact structures

Eliashberg (1989): fundamentally 2 types of contact structures.

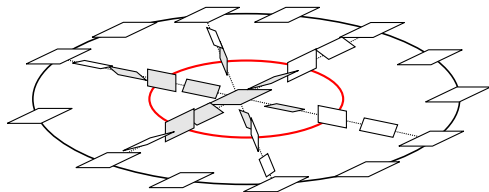
- *Overtwisted*: contains an *overtwisted disc*.
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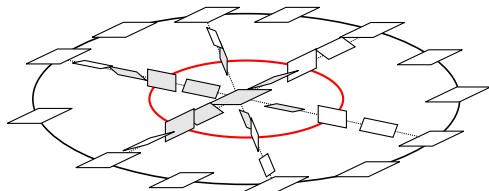


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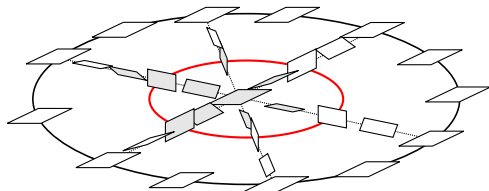
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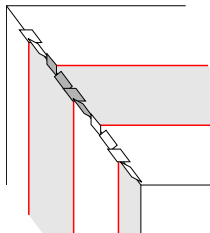
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- Overtwisted contact geometry reduces to (well-understood) homotopy theory. Tight contact structures offer important topological information.
- Eliashberg (1992): contact structure near an S^2 is tight iff dividing set is *connected*. If so, contact structure extends uniquely (up to isotopy) to a tight contact structure on B^3 .

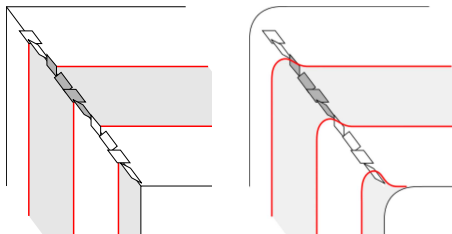
Contact corners

When two convex surfaces meet along a boundary, contact planes are arranged as shown.



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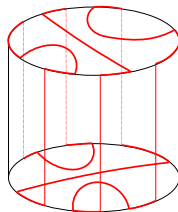
Proposition

Let Γ_0, Γ_1 be chord diagrams. The following are equivalent:

- $\langle \Gamma_0 | \Gamma_1 \rangle = 1$.
- The solid cylinder with dividing set Γ_0 on the bottom and Γ_1 on the top has a tight contact structure.

Bypasses

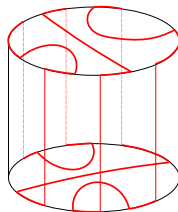
Honda (2000's): any 3-manifold can be built up from a surface and dividing set by adding *bypasses*.



Effect on dividing set is “bypass surgery” as defined earlier.

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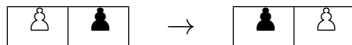
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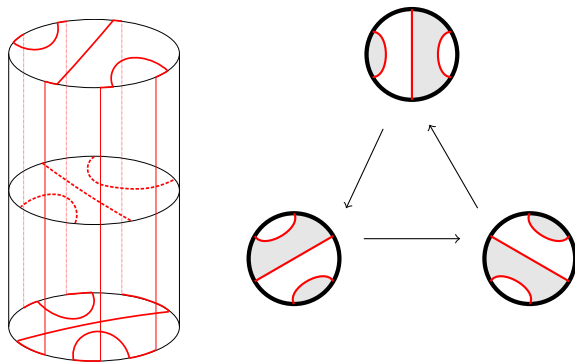
$$\langle \Gamma_{pq} | \Gamma_{qp} \rangle = 1$$

or



Bypasses

Stacking two bypasses on top of each other produces an overtwisted contact structure!



Can build something like a *triangulated category* out of dividing sets and contact structures (Honda, M.). V_n is the *Grothiendick group*.

Contact TQFT = Quantum pawn dynamics

These definitions give many of the properties of a (2+1)-dimensional *topological quantum field theory*.

- Contact structure near disc (2-dim) \rightsquigarrow “states” in V_n
- Contact structure over cylinder (2+1-dim) \rightsquigarrow element of \mathbb{Z}_2 .
- “Possibility of a tight contact structure from one state to another” \rightsquigarrow inner product $\langle \cdot | \cdot \rangle : V_n \otimes V_n \longrightarrow \mathbb{Z}_2$.

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Theorem (M.)

“Contact TQFT is isomorphic to quantum pawn dynamics.”

Thanks for listening!

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